

What is a Computer?

Programming, Memory and Limits of Computation

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10.22.13

Outline

Introduction

Finite State Machines

Computers

Programming

Computation and Turing Machines

Extents and Limits of Computation

Outline

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Finite State Machines

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Extents and Limits of Computation

Review

From Minds to Machines

- Dualism 🖐️, Identity Theory 🖐️, Functionalism 👍

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- ▶ So what exactly is computation?
- ▶ Could the things minds do amount to computation?

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- ▶ And FSMs are another important building block of computers

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Finite State Machines

Computers

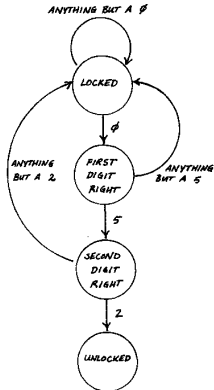
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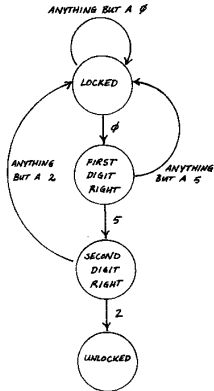
Finite State Machines

Combo Lock For 0-5-2 (p.36)



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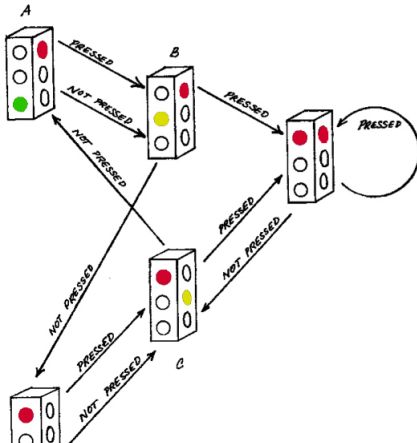


FSMs Listen to Themselves

- ▶ Response = input + **current state**
- ▶ Current state = **previous input**

Stoplight Controller

As a Finite State Machine



How A Stoplight Works

In Terms of Input/Output

Inputs:

**Walk
Button**

Not Pressed
Not Pressed
Not Pressed
Not Pressed
Not Pressed
Pressed
Pressed
Pressed
Pressed
Pressed

**Current
State**

A
B
C
D
Walk
A
B
C
D
Walk

**Main
Road**

Red
Red
Yellow
Green
Walk
Red
Red
Yellow
Green
Walk

Outputs:

**Cross
Road**

Green
Yellow
Red
Red
Walk
Green
Yellow
Red
Red
Walk

**Next
State**

B
D
A
C
D
B
Walk
Walk
C
Walk

Finite State Machines

The Mealy Machine

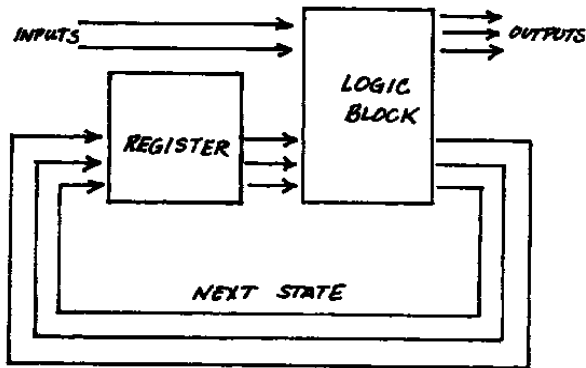


FIGURE 14

Finite State Machines

In General

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1. **Register** built from Boolean circuits

Finite State Machines

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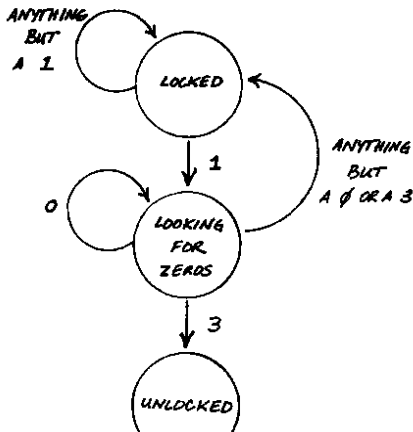
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 - ▶ All patterns?

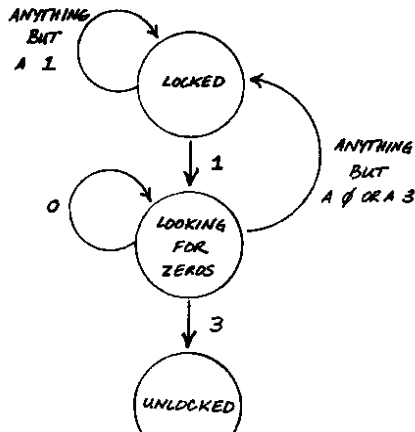
Finite State Machines

Recognizing Patterns: begins w/1, any number of 0s, ends w/3 (p.35)



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Important

There is one state for each digit that has to be remembered.

Minds Aren't Just Logic Blocks

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- ▶ The register itself is a logic block: a configuration of bits that changes upon input

Finite State Machines

Recognizing Patterns

Palindrome

A string of symbols which is the same whether read backwards or forwards.

- ▶ E.g. 99, 757, 1001, abba, abccba, dad, mom

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- ▶ A finite state combo lock that detects only palindromes is not possible (even if alphabet is finite)

Finite State Machines

Recognizing Patterns

Why

To recognize a palindrome, you need to remember **every** digit of the first half. You need one state for each such memory. But there are infinitely many first halves, so you would need infinitely many states.

Finite State Machines

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- ▶ Isn't this a sentence: *I don't believe that S ?*

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From Finite State Machines to Computers

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- ▶ Final ingredient of a computer: a more sophisticated kind of **memory**

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Finite State Machines

In General

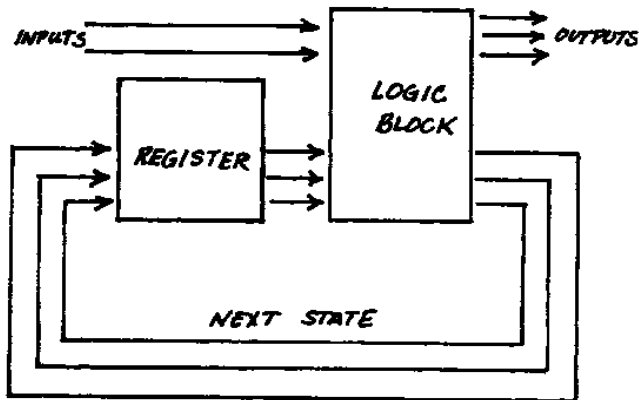
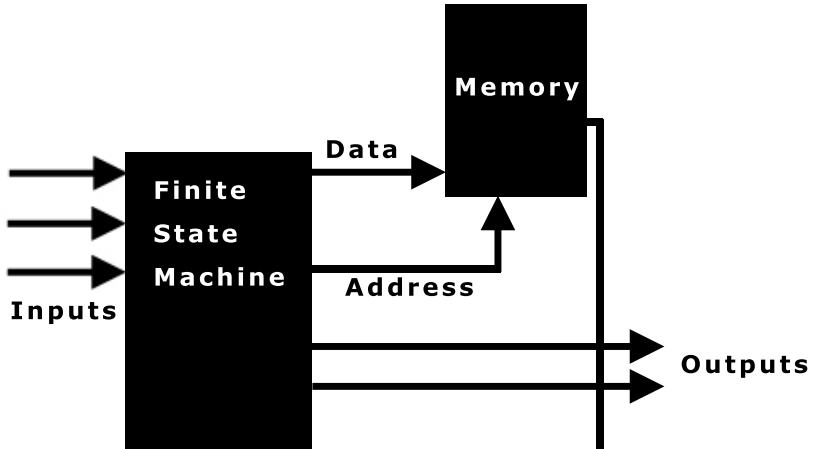


FIGURE 1.1

Finite State Machine

With a Real Memory



What is A Real Memory?

Memory, Registers, Words, Addresses, Reading, Writing

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- ▶ This bit pattern is either an instruction or some data
- ▶ Data are treated as input, instructions are executed

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- ▶ For instance, the FSM may react to a particular bit pattern from memory by adding
- ▶ This involves first getting bit patterns from inputs or registers
- ▶ A logic block then adds these bit patterns, and writes the result to the memory

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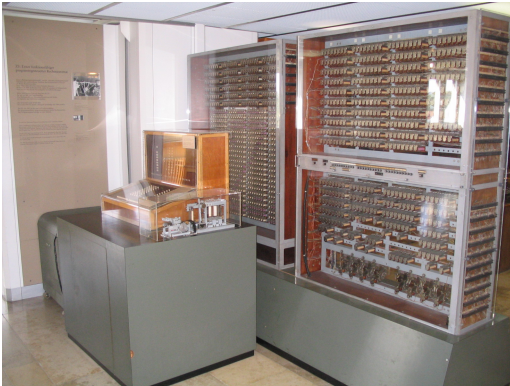
Programming

Computation and Turing Machines

Extents and Limits of Computation

A General Purpose Computer

How Do You Get a Computer To Do Your Bidding?



- ▶ How does a computer interpret bit patterns that are to be treated as instructions?
- ▶ How do you program one to do new things?

Zuse Z3 (1941): 1st Programmable Computer

Machine Language

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- ▶ Responses to **processing instructions** are: move data to and from memory, perform arithmetic, etc.
- ▶ Responses to **control instructions** determine address of next instruction to be retrieved
 - ▶ Stored in special register: **program counter**
- ▶ Normally computer will just move from one register to next, but control instructions can tell it to skip some or repeat some until a condition is met (loop)

Programming Languages

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- ▶ So how do you teach your computer to do new things like record music or play video's?

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- ▶ The computer does this by translating the programming language into its machine language

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Logo and the Tortoise

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 - ▶ And lists next to them the addresses of the machine code commands the Logo commands correspond to
- ▶ When a sequence corresponding to a Logo command is entered, the FSM looks up and performs the corresponding commands of machine code
- ▶ So Logo is just a convenient Macro!

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6. The FSM and memory are built from logic blocks
7. Logic blocks crunch bits
8. Crunching bits is manipulating a physical substance to send one of two possible signals 1 or 0

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Today

Overview

1. Learn Turing's way of thinking about the memory

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5. Learn about the limits of **computability**:
 - ▶ **Halting Problem**: no computer will ever solve it

Calculation

A Person, A Pencil and Some Paper

- You start by writing some things down

$$X - Z = Y$$



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- ▶ You start by writing some things down
- ▶ You follow some rules for using them to get some thing else
- ▶ You do this by writing out each application of the rule
- ▶ These rules have the form: if you've written X , you can write Y

$$X \rightarrow Z = Y$$



Calculation and Computation

Turing's Staring Point

Effective Calculation

An **effective calculation** is a calculation that:

1. Is set out in a finite number of exact instructions

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Calculation and Computation

Turing's Starting Point

Effective Calculation

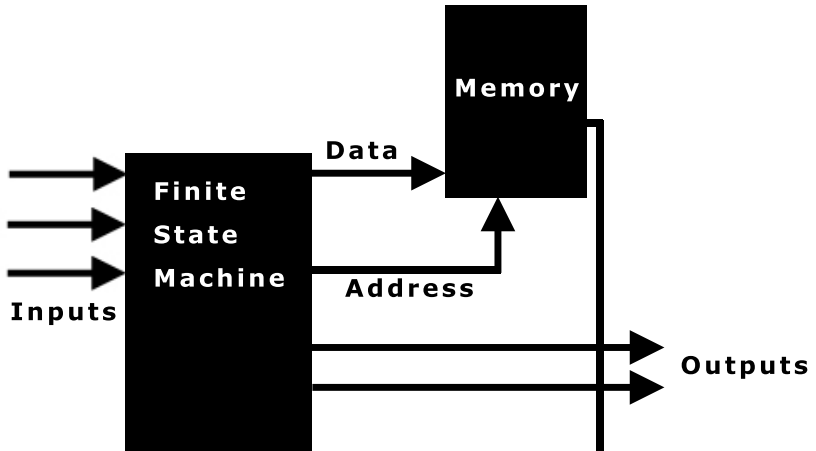
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Turing: to compute is just to effectively calculate

A Computer

FSM with a Memory



Computation and Calculation

Turing Machines

- Is a computer really doing (effective) calculations?

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- ▶ The **Turing Machine**:
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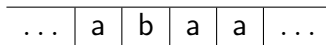
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| | | | | | |
|-----|---|---|---|---|-----|
| ... | a | b | a | a | ... |
|-----|---|---|---|---|-----|

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- ▶ 3 **machine commands**: write, move head left (\ll), move head right (\gg)

Turing Machines

And Programs

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Turing Machines

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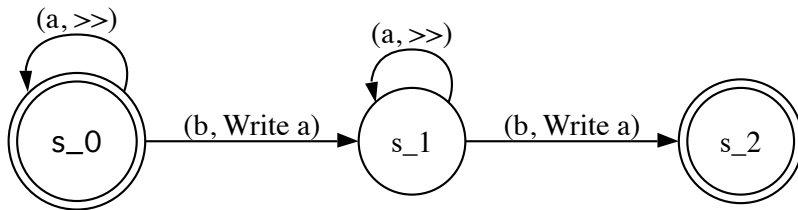
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| | a | b |
|-------|------------|----------------|
| s_0 | \gg, s_0 | Write a, s_1 |
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Visualizing A Turing Machine

Executing This Program

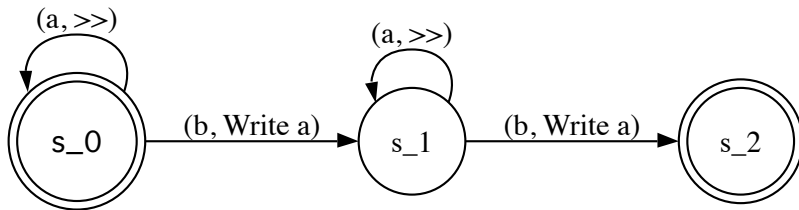
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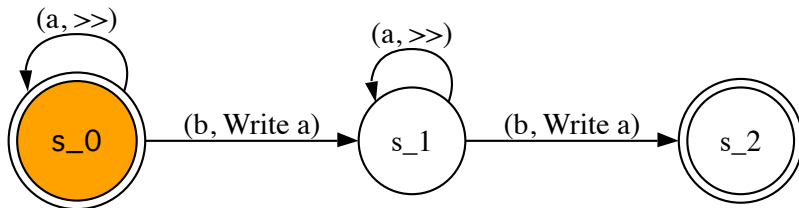
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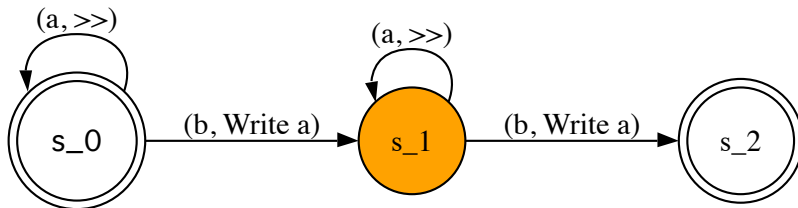
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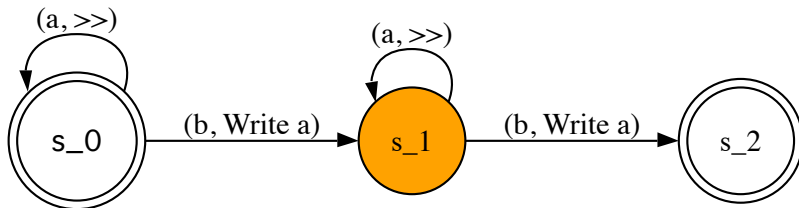
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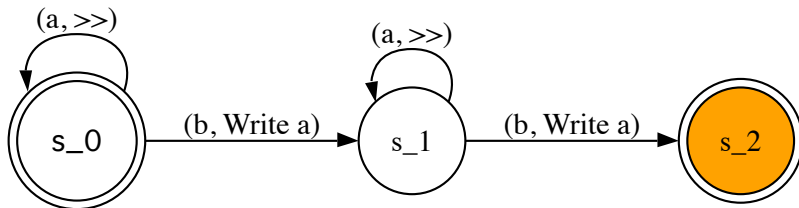
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When a Turing Machine executes a program, what part of the machine is causing it to follow the routine? Where is the program in the machine and how does it cause the memory and states to change?

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- ▶ A Turing Machine can be **programmed** by constructing the **internal configuration** and **wiring** of the **FSM** in a particular way

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- ▶ Recall that the Turing Machine consists of the tape assembly wired to a FSM
- ▶ A Turing Machine can be **programmed** by constructing the **internal configuration** and **wiring** of the **FSM** in a particular way
- ▶ But Turing showed there's **another possibility**

Programs and The Universal Turing Machine

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Universal Turing Machine (UTM)

Given the appropriate tape, the UTM can do anything that any other Turing Machine can do. With the right software, it can do anything! (Though maybe much slower!)

The Church-Turing Thesis

Effective Calculation

The Church-Turing Thesis (Widely Accepted Conjecture)

Any effective calculation can be performed by a Turing Machine.

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Outline

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Finite State Machines

Computers

Programming

Computation and Turing Machines

Extents and Limits of Computation

Alternative Computation?

Is Turing All There Is?

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- ▶ Well, think back to our discussion of Rock-Paper-Scissors

Rock-Paper-Scissors

Not A Binary Input, Not A Binary Output!

| Input A | Input B | Output |
|----------------|----------------|---------------|
| Scissors | Scissors | Tie |
| Scissors | Paper | A wins |
| Scissors | Rock | B wins |
| Paper | Scissors | B wins |
| Paper | Paper | Tie |
| Paper | Rock | A wins |
| Rock | Scissors | A wins |
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| Rock | Rock | Tie |

Rock-Paper-Scissors

Using Two Binary Inputs to Mimic One Ternary Input

| | A Inputs | B Inputs | Outputs |
|----------------------|----------|----------|---------|
| | 01 | 01 | 00 |
| | 01 | 10 | 10 |
| Scissors = 01 | 01 | 11 | 01 |
| Paper = 10 | 10 | 01 | 01 |
| Rock = 11 | 10 | 10 | 00 |
| A wins = 10 | 10 | 11 | 10 |
| B wins = 01 | 11 | 01 | 10 |
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- ▶ These exist and are called **analog computers**
- ▶ So are they more powerful?

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Noise

- ▶ Instead of responding differently to some/no voltage, it will respond differently to any fraction of voltage differently

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- ▶ To control for this noise you have to ignore some of the nuances of the signal

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- ▶ Could computers somehow be built to harness the randomness of the quantum, sub-atomic world?
 - ▶ The jury's still out on that one

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Something Computers Can't Do

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Is it possible to write a program that inspects any other programs and determines whether or not they will eventually stop?

- ▶ Suppose it is, call it Test-for-Halt

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- ▶ Suppose Test-for-Halt says 'yes'
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The Halting Problem

Something Computers Can't Do

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Answer: No, it's not possible. You can always build a

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- ▶ An algorithm is a fail-safe procedure that, if executed correctly, guarantees the result.
- ▶ Programs are representations of algorithms in some programming language.
- ▶ There are usually several ways of expressing in algorithms, depending on the programming language.

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Answer: We can describe the speed of an algorithm on the basis of how much the time required for performing the task grows along with the size of the problem.

Algorithms

Comparing Algorithms

First sock-matching algorithm

- ▶ Pull out a random sock out of the basket.
- ▶ Pull out a second sock out of the basket.
- ▶ Compare it with the first pulled out.
- ▶ If it matches, bundle them together.
- ▶ if it does not match, throw it back in the basket and restart the process.

Algorithms

Comparing Algorithms

Second sock-matching algorithm

- ▶ Pull out a random sock out of the basket on the table.
- ▶ Pull out a second sock out of the basket.
- ▶ Compare it with the first pulled out.
- ▶ If it matches with any sock on the table, bundle them together.
- ▶ if it does not match, set it next to the first.
- ▶ Repeat the process until all the socks are matched.

Algorithms

Comparing Algorithms

Compare the first and the second algorithm

- ▶ Assume that there are n socks in the basket.
- ▶ Using the first algorithm, finding two that match requires pulling out and pulling back an average of half of the remaining socks.
- ▶ So the number of sock removals is proportional to the square of number of socks.
- ▶ The algorithm is of order n^2 , meaning that for large problem, the time of execution grows as the square of the problem size.

Algorithms

Comparing Algorithms

Compare the first and the second algorithm

- ▶ In other words, if there are m times as many socks in the basket, the time of execution of the second algorithm will only by m by n .
- ▶ The time of execution of the first algorithm will be instead m^2 by n .
- ▶ So if $n=10$, then using the second algorithm to match twice as many socks will take twice as much time.
- ▶ Instead, if $n=10$, then using the first algorithm to match twice as many socks will take four times as much time.

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